

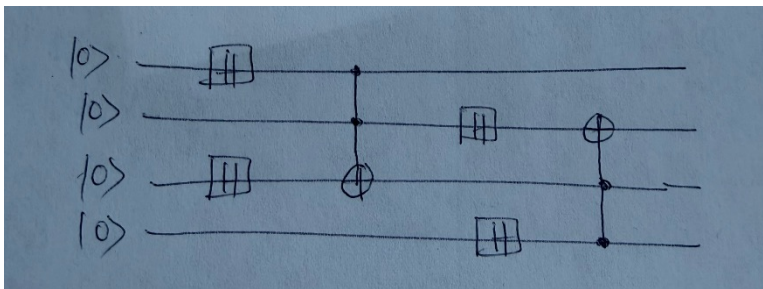
Name: \_\_\_\_\_

## CSE 4502/5717 Big Data Analytics

### Exam III (model); Spring 2026

**Note:** You are supposed to give proofs to the time and processor bounds of your algorithms. Read the questions carefully before attempting to solve them.

1. Input is a database DB with  $n$  transactions from a set  $I = \{i_1, i_2, \dots, i_d\}$  of items. Input also is a threshold  $\text{minSupport}$  for the minimum support. It is known that each item in  $I$  occurs in at most  $c$  transactions, where  $c$  is a constant. Present an  $O(d)$  time algorithm to identify all the frequent items in DB.
2. Let DB be a database with  $n$  transactions. Let  $d$  be the number of possible items.  $X$  is a specific  $k$ -itemset. The problem is to compute an estimate on  $q = \sigma(X)$ , i.e., the number of transactions in which  $X$  occurs. It is known that  $\sigma(X) \geq \sqrt{n}$ . Present an  $O(k\sqrt{n} \log n)$  time algorithm to come up with an estimate  $\mu$  such that  $0.9q \leq \mu \leq 1.1q$ . Show that  $\mu$  will have this property with a high probability (i.e., a probability of  $\geq (1 - n^{-\alpha})$ , for any fixed  $\alpha$ ). Assume that each transaction is given as a bit array as discussed in class.
3. Input is a database DB with  $n$  transactions from a set  $I = \{i_1, i_2, \dots, i_d\}$  of items. Input also is a threshold  $\text{minSupport}$  for the minimum support. We are required to identify all the frequent  $k$ -itemsets, where  $k$  is a constant. Present a parallel algorithm for this problem that runs in  $O(\log n)$  time. You can use up to  $\frac{nd^k}{\log n}$  CREW PRAM processors. Assume that each transaction is given as a bit array.
4. Show how to solve the hierarchical clustering problem in  $O(n \log n)$  time using  $\frac{n}{\log n}$  CREW PRAM processors.
5. What is the output of the following quantum circuit:



6. Input is a not necessarily sorted sequence  $X$  of  $n$  arbitrary elements. The problem is to find the minimum of  $X$ . Show that this problem can be solved using Grover's algorithm in  $O(\sqrt{n})$  time with a constant probability. Given a sequence  $X = x_1, x_2, \dots, x_N$  and a function  $f : X \rightarrow \{0, 1\}$ , Grover's algorithm finds an  $i$  such that  $f(x_i) = 1$ , in  $O(\sqrt{N})$  time. If there are  $k$  values for  $i$  such that  $f(x_i) = 1$ , assume that Grover's algorithm runs in time  $O(\sqrt{N/k})$ .